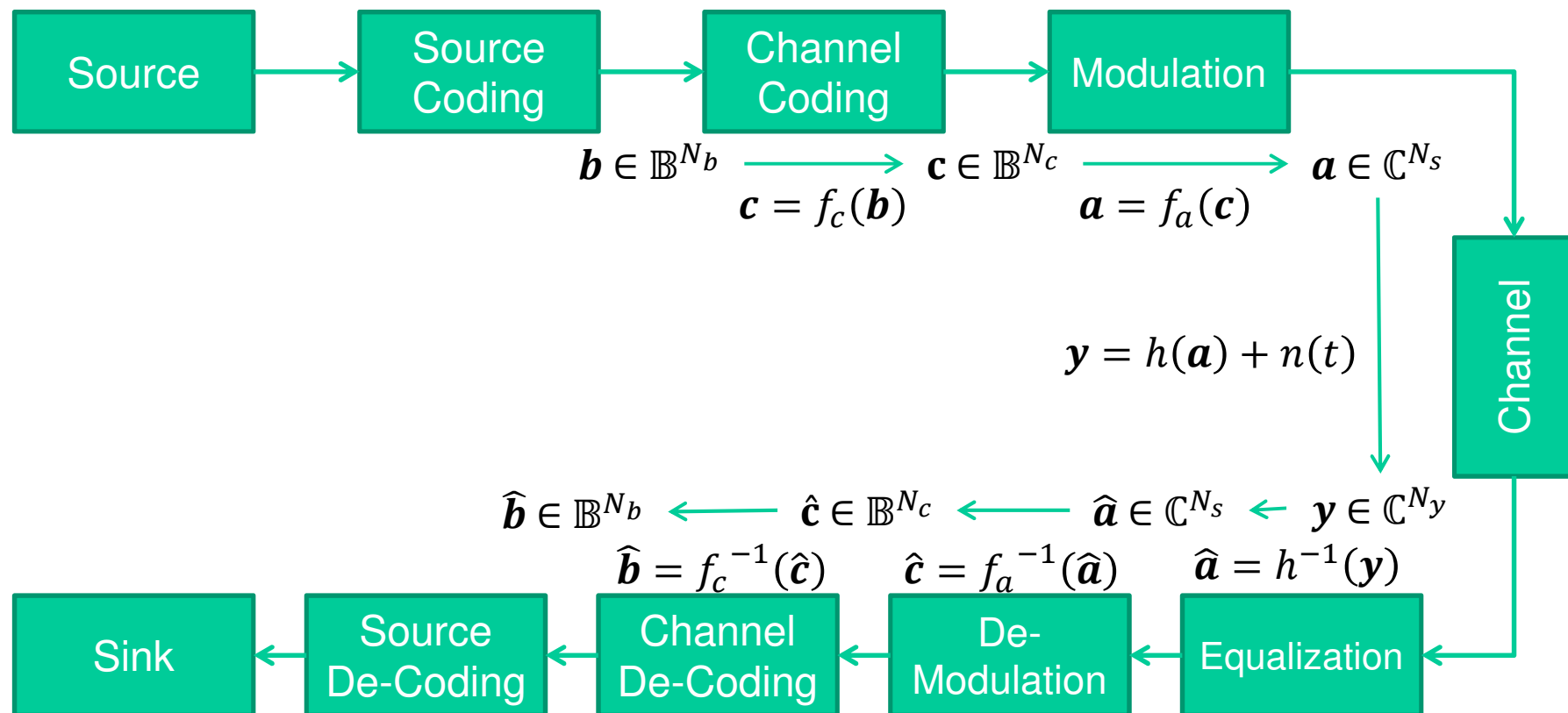


Iterative Receiver

Outline

Digital Communication
From Receiver Design to Factor Graph
Channel-Decoding
De-Modulation
Equalization
Synchronization
All-In-One
Summary

Classical Digital Communication



Baseband Equivalent Model

Receiver Estimation Process

Aim of a receiver: Estimation of the transmitted data (bits) in an “optimal” sense

- Sequence/Frame Detection

$$\hat{\mathbf{b}} = \arg \max_{\mathbf{b} \in \mathbb{B}^{N_b}} \{p(\mathbf{B} = \mathbf{b} | \mathbf{Y} = \mathbf{y}, \mathcal{M})\}$$

- Maximizing the a-posteriori probability of a data frame is equivalent to minimizing the frame error rate (FER)

- Bit-by-Bit Detection

$$\hat{b}_k = \arg \max_{b_k \in \mathbb{B}} \{p(B_k = b_k | \mathbf{Y} = \mathbf{y}, \mathcal{M})\}$$

- Maximizing the a-posteriori probability of a single data bit is equivalent to minimizing the bit error rate (BER)

- →MAP Estimation

Factor Graphs in Digital Communication

- The needed A-Posteriori $p(B_k = b_k | Y = y, \mathcal{M})$ is given as:

$$p(B_k = b_k | Y = y, \mathcal{M}) = \frac{p(B_k = b_k, Y = y | \mathcal{M})}{\sum_b p(B_k = b_k, Y = y | \mathcal{M})}$$

- Factorization of the joint Probability $p(\mathbf{B}, Y = y | \mathcal{M})$:

$$p(\mathbf{B}, Y = y | \mathcal{M}) = p(Y = y | \mathbf{B}, \mathcal{M}) p(\mathbf{B} | \mathcal{M})$$

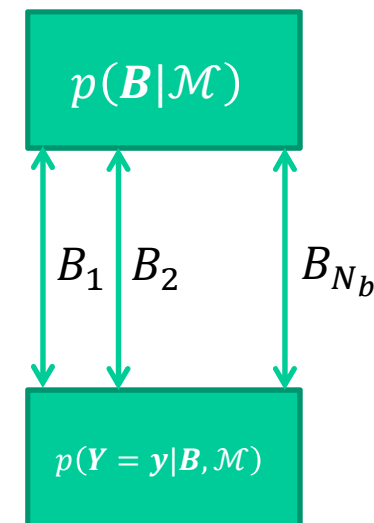
- A-Priori knowledge of the transmitted Bits:

$$p(\mathbf{B} = \mathbf{b} | \mathcal{M}) = \prod_{k=1}^{N_b} p(B_k = b_k | \mathcal{M})$$

- How does the Likelihood $p(Y = y | \mathbf{B}, \mathcal{M})$ look like:

$$p(Y = y | \mathbf{B}, \mathcal{M}) = ???$$

→ “Opening” the Factor!



Channel De-Coding (1)

- „Opening“ the Factor $p(Y = \mathbf{y}|\mathbf{B}, \mathcal{M})$:

$$p(Y = \mathbf{y}|\mathbf{B}, \mathcal{M}) = \sum_{\mathbf{c}} p(Y = \mathbf{y}, \mathbf{C} = \mathbf{c}|\mathbf{B}, \mathcal{M})$$

- Factorization of the joint Probability $p(Y = \mathbf{y}, \mathbf{C}|\mathbf{B}, \mathcal{M})$:

$$\begin{aligned} p(Y = \mathbf{y}, \mathbf{C}|\mathbf{B}, \mathcal{M}) &= p(Y = \mathbf{y}|\mathbf{C}, \mathbf{B}, \mathcal{M})p(\mathbf{C}|\mathbf{B}, \mathcal{M}) \\ &= p(Y = \mathbf{y}|\mathbf{C}, \mathcal{M})p(\mathbf{C}|\mathbf{B}, \mathcal{M}) \end{aligned}$$

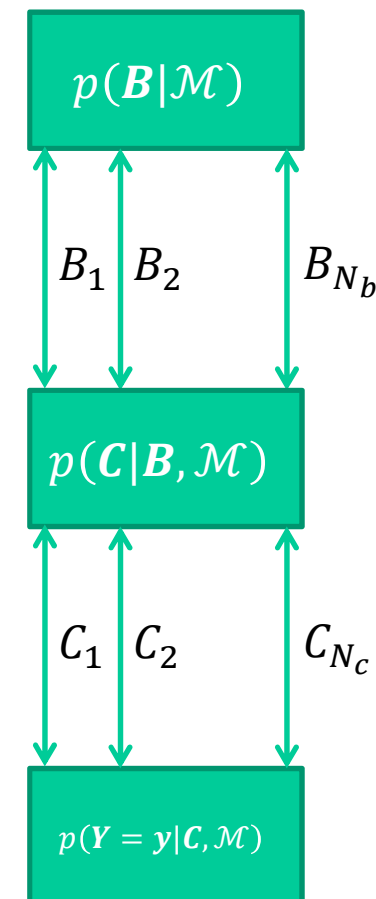
- The factor $p(\mathbf{C}|\mathbf{B}, \mathcal{M})$ describes the coding scheme:

$$p(\mathbf{C} = \mathbf{c}|\mathbf{B} = \mathbf{b}, \mathcal{M}) = \mathbb{I}(\mathbf{c}, f_c(\mathbf{b}))$$

$$\mathbb{I}(\mathbf{c}, f_c(\mathbf{b})) = \begin{cases} 1, & \mathbf{c} = f_c(\mathbf{b}) \\ 0, & \mathbf{c} \neq f_c(\mathbf{b}) \end{cases}$$

- Again: How does the Likelihood $p(Y = \mathbf{y}|\mathbf{C}, \mathcal{M})$ look like:

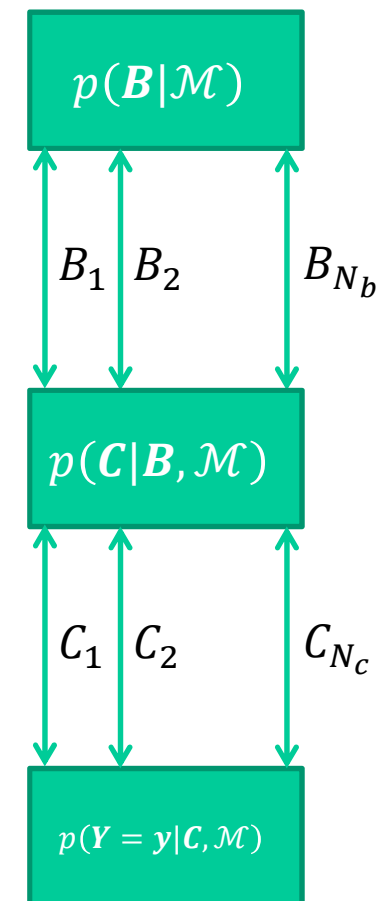
$$p(Y = \mathbf{y}|\mathbf{C}, \mathcal{M}) = ???$$



Channel De-Coding (2)

Example (1/4 Repetition Code):

- Known:
 - Data Prior:
$$p(B = 0) = 0.5$$
 - Channel Likelihood:
$$p(Y = 0|C = 0) = 0.5$$
$$p(Y = 0|C = 1) = 0.1$$
- Given Observation:
$$\mathbf{y} = [0, 0, 1, 1]$$
- Find: \mathbf{b}

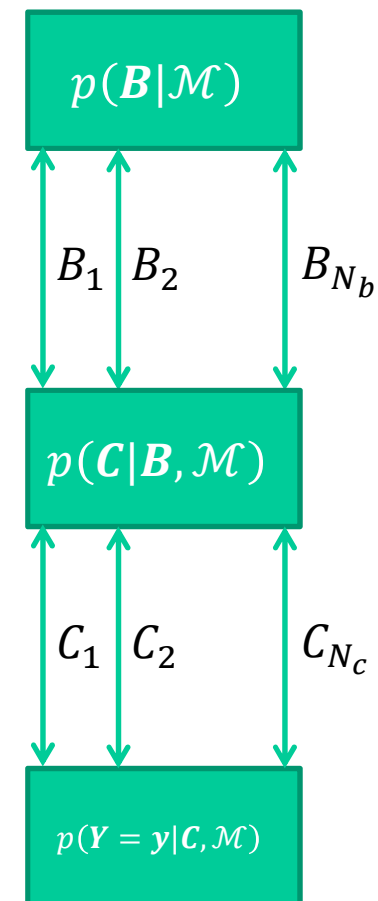


Channel De-Coding (3)

Example (1/4 Repetition Code),
ctd:

$$\begin{aligned}
 p(B, Y = \mathbf{y}) &= p(Y = \mathbf{y} | B) p(B) \\
 &= \sum_{\mathbf{c}} p(Y = \mathbf{y}, \mathbf{C} = \mathbf{c} | B) p(B)
 \end{aligned}$$

$$\begin{aligned}
 p(Y = \mathbf{y}, \mathbf{C} | B) &= p(Y = \mathbf{y} | \mathbf{C}, B) p(\mathbf{C} | B) \\
 &= \prod_k p(Y_k = y_k | C_k) \mathbb{I}(C_k, B)
 \end{aligned}$$



Channel De-Coding (4)

Example (1/4 Repetition Code),
ctd:

with $\mathbf{y} = [0,0,1,1]$

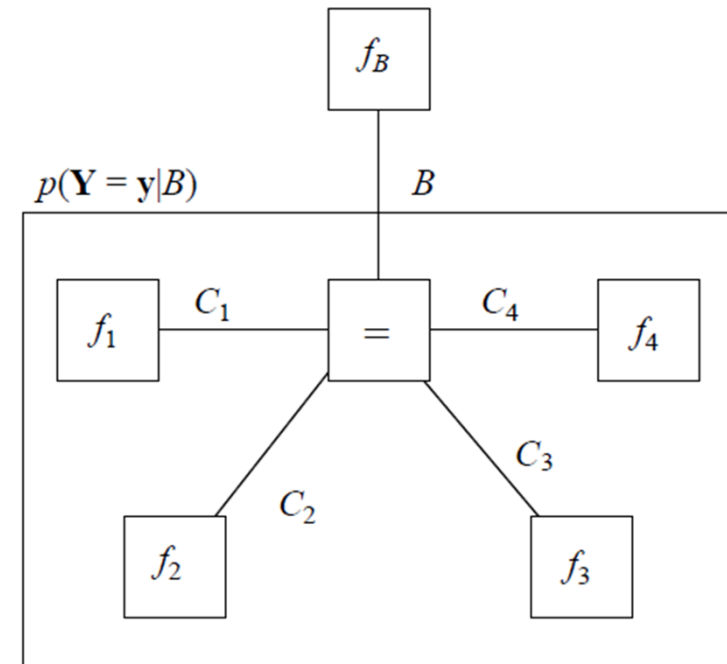
$$\mu_{f_B \rightarrow B}(b) = \begin{cases} 0.5; & b = 1 \\ 0.5; & b = 0 \end{cases}$$

for $k = 1, 2$

$$\mu_{f_k \rightarrow B}(b) = \begin{cases} 0.1; & b = 1 \\ 0.5; & b = 0 \end{cases}$$

for $k = 3, 4$

$$\mu_{f_k \rightarrow B}(b) = \begin{cases} 0.9; & b = 1 \\ 0.5; & b = 0 \end{cases}$$



Channel De-Coding (5)

Example (1/4 Repetition Code), ctd:

$$\mu_{\Rightarrow B}(B = b) = \sum_c \mathbb{I}(b, c_1, c_2, c_3, c_4) \prod_k \mu_{f_k \rightarrow C_k}(c_k) = \prod_k \mu_{f_k \rightarrow C_k}(b)$$

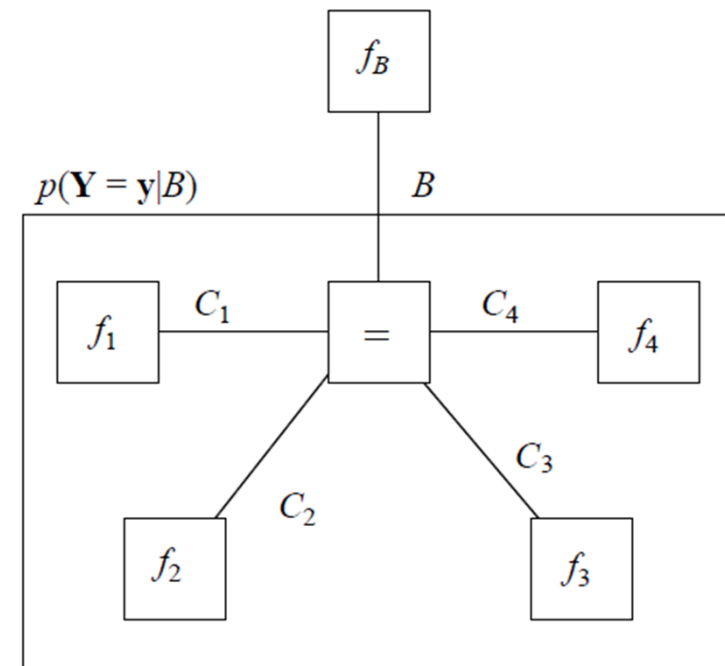
$$\begin{aligned} \mu_{\Rightarrow B}(B = 0) &= (0.5)^4 = 0.0625 \\ \mu_{\Rightarrow B}(B = 1) &= (0.9)^2(0.1)^2 = 0.0081 \end{aligned}$$

$$g_B(B = b) = \mu_{f_B \rightarrow B}(b) \mu_{\Rightarrow B}(b) = p(B = b, \mathbf{Y} = \mathbf{y}):$$

$$\begin{aligned} g_B(B = 1) &= 0.00405 \\ g_B(B = 0) &= 0.03125 \end{aligned}$$

$$p(B = b | \mathbf{Y} = \mathbf{y}) = \frac{p(B = b, \mathbf{Y} = \mathbf{y})}{\sum_b p(B = b, \mathbf{Y} = \mathbf{y})}:$$

$$\begin{aligned} p(B = 1 | \mathbf{Y} = \mathbf{y}) &= \mathbf{0.11} \\ p(B = 0 | \mathbf{Y} = \mathbf{y}) &= \mathbf{0.89} \end{aligned}$$



De-Modulation (1)

- „Opening“ the Factor $p(Y = \mathbf{y}|\mathbf{C}, \mathcal{M})$:

$$p(Y = \mathbf{y}|\mathbf{C}, \mathcal{M}) = \sum_{\mathbf{c}} p(Y = \mathbf{y}, \mathbf{A} = \mathbf{a}|\mathbf{C}, \mathcal{M})$$

- Factorization of the joint Probability $p(Y = \mathbf{y}, \mathbf{A}|\mathbf{C}, \mathcal{M})$:

$$\begin{aligned} p(Y = \mathbf{y}, \mathbf{A}|\mathbf{C}, \mathcal{M}) &= p(Y = \mathbf{y}|\mathbf{A}, \mathbf{C}, \mathcal{M})p(\mathbf{A}|\mathbf{C}, \mathcal{M}) \\ &= p(Y = \mathbf{y}|\mathbf{A}, \mathcal{M})p(\mathbf{A}|\mathbf{C}, \mathcal{M}) \end{aligned}$$

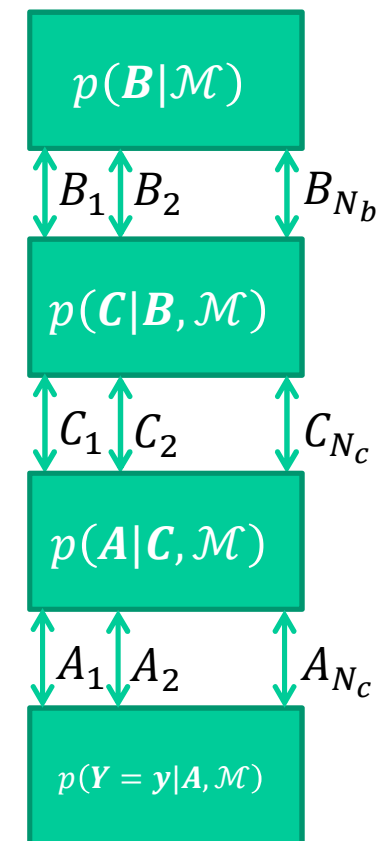
- The function $p(\mathbf{A}|\mathbf{C}, \mathcal{M})$ describes the modulation scheme:

$$p(\mathbf{A} = \mathbf{a}|\mathbf{C} = \mathbf{c}, \mathcal{M}) = \mathbb{I}(\mathbf{a}, f_a(\mathbf{c}))$$

$$\mathbb{I}(\mathbf{a}, f_a(\mathbf{c})) = \begin{cases} 1, & \mathbf{a} = f_a(\mathbf{c}) \\ 0, & \mathbf{a} \neq f_a(\mathbf{c}) \end{cases}$$

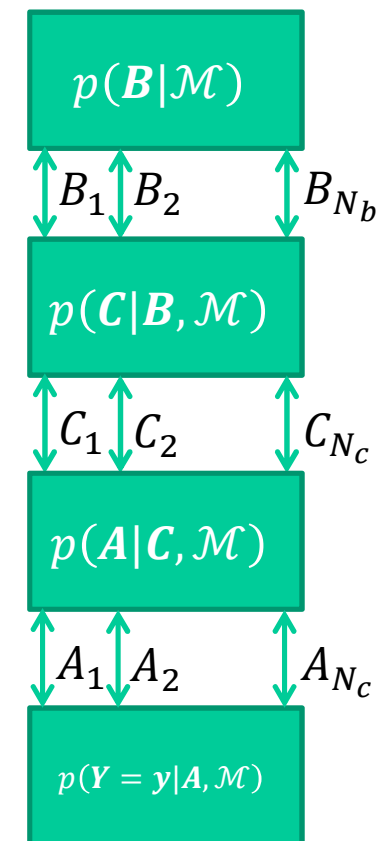
- And Again: How does the Likelihood $p(Y = \mathbf{y}|\mathbf{A}, \mathcal{M})$ look like:

$$p(Y = \mathbf{y}|\mathbf{A}, \mathcal{M}) = ???$$



Equalization (1)

- The likelihood function $p(Y = \mathbf{y} | \mathbf{A}, \mathcal{M})$ gives the likelihood that the observed signal \mathbf{y} represents a given modulation symbol \mathbf{a} .
- Equalization estimate the channel which has an influence on the transmitted symbol



Equalization (2)

Example (BPSK):

BPSK symbol $a \in \{-1, +1\}$ with the mapping function

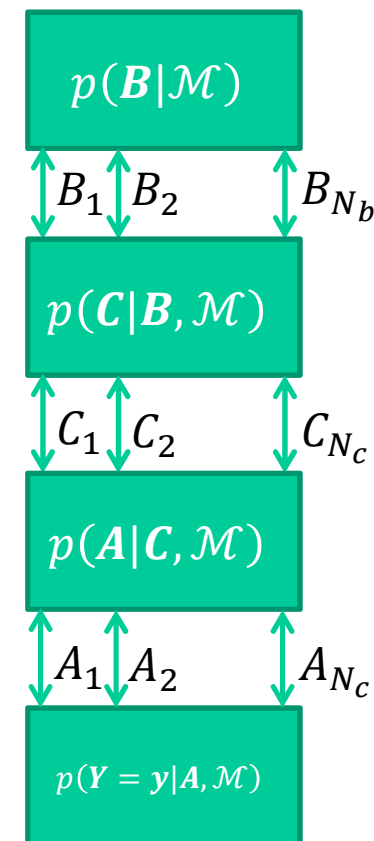
$$a_n = f_c(c_n) = 2c_n - 1$$

Considering AWGN channel

$$\mathbf{y} = \mathbf{a} + \mathbf{n}$$

with \mathbf{n} is a RV with $\mathcal{N}(0, \sigma^2)$:

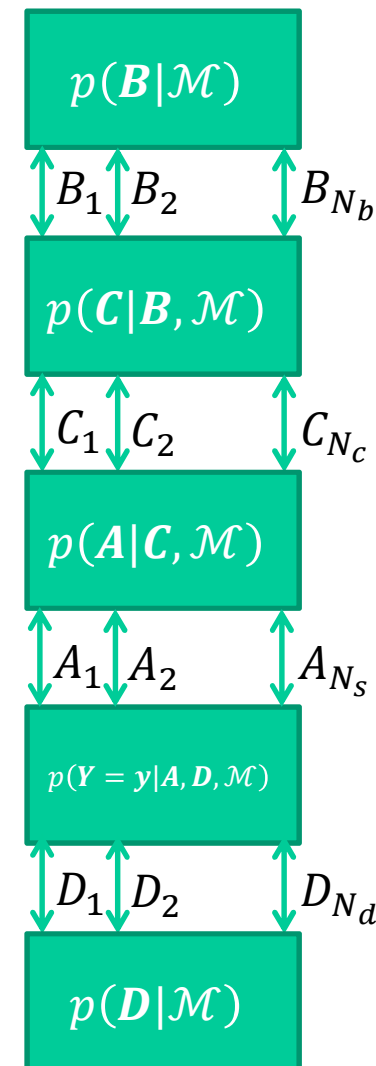
$$p(Y_k = y_k | A_k, \mathcal{M}) \propto \exp\left(-\frac{(y_k - A_k)^2}{2\sigma^2}\right)$$



Synchronization (1)

- Until now the (deterministic) channel h was considered to be known
- Synchronization introduces an statistical channel approach using the channel parameters \mathbf{D}
- Therefore the factor $p(\mathbf{Y} = \mathbf{y}|\mathbf{A}, \mathcal{M})$ is expanded for due to the statistical channel model

$$p(\mathbf{Y} = \mathbf{y}|\mathbf{A}, \mathbf{D}, \mathcal{M}) = p(\mathbf{Y} = \mathbf{y}|\mathbf{A}, \mathcal{M}) p(\mathbf{D}|\mathcal{M})$$



Synchronization (2)

Example (Multipath Channel):

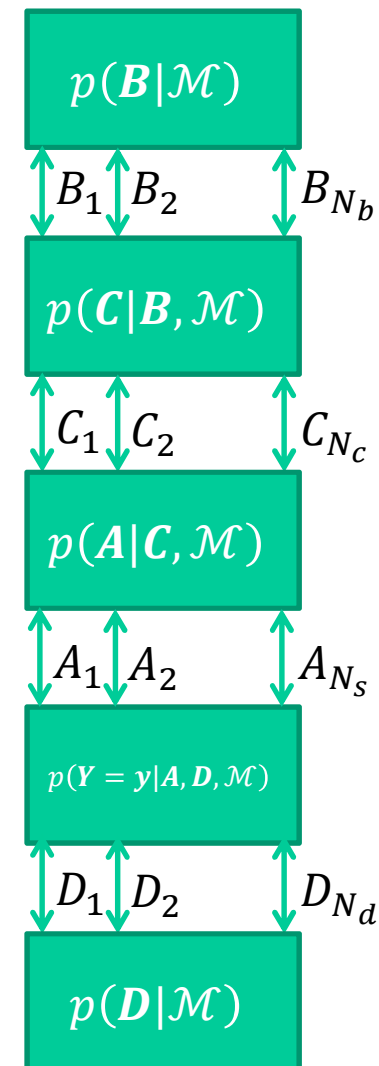
- The multipath channel is typically denoted as

$$h(t) = \sum_{k=0}^K \alpha_k \delta(t - \tau_k)$$

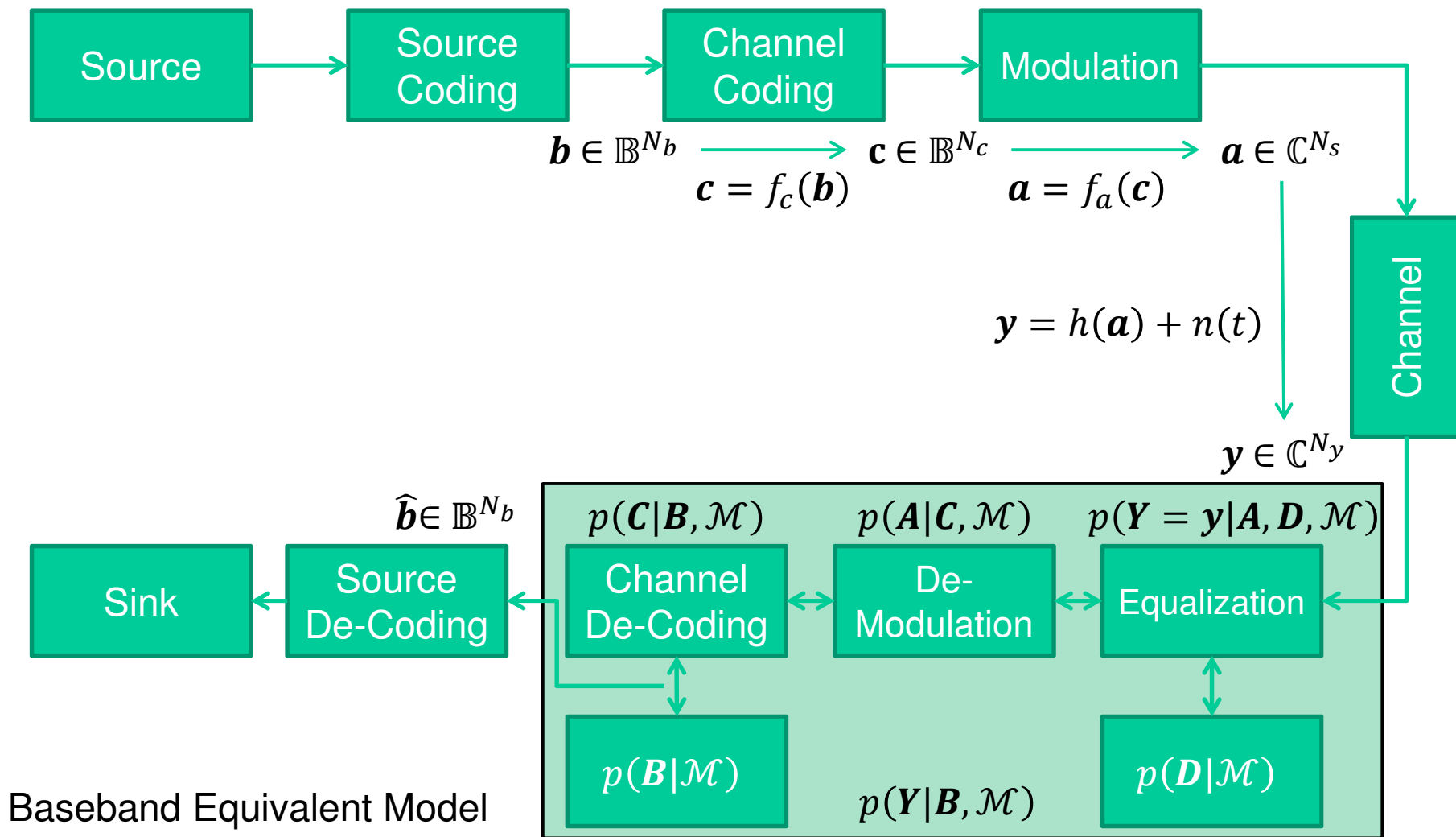
- Therefore:

$$\mathbf{D} \in \{\boldsymbol{\alpha}, \boldsymbol{\tau}\}^{K+1}$$

with $\boldsymbol{\alpha}, \boldsymbol{\tau}$ are RV with specific distributions



Iterative Receiver Design



Summary

- A classical receiver design can be fully substituted by an iterative approach
- Factor Graphs provide a very good framework to design the different stages e.g.
 - Source De-Coding
 - De-Modulation
 - Equalization
 - Synchronization
- The full receiver factor graph has loops (also between different blocks)
- Iterative Receiver Design gives the ability to distribute estimation information „backward“